

SEMIDEFINITE APPROXIMATIONS FOR BIINDEPENDENT PAIRS AND BICLIQUES IN BIPARTITE GRAPHS

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Joint work with **Sven Polak** (Tilburg University)
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Biindependent pairs in a bipartite graph

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Let $G = (V_1 \cup V_2, E)$ be a **bipartite** graph. A pair (A, B) of subsets $A \subseteq V_1, B \subseteq V_2$ is **biindependent** if $\{i, j\} \notin E \quad \forall (i, j) \in A \times B$.

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A pair (A, B) is **balanced** if $|A| = |B|$.

Define the **balanced parameters** $\alpha_{\text{bal}}(G)$, $g_{\text{bal}}(G)$, $h_{\text{bal}}(G)$,
where we optimize over **balanced biindependent** pairs.

Extension to bicliques and to general graphs

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Biindependent pairs in bipartite graphs permit to model biindependent pairs and bicliques in general graphs

Definition (biindependent pairs and bicliques in general graphs)

Let $G = (V, E)$ be a (general) graph and let $A, B \subseteq V$ be disjoint.

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Lemma (Links to bipartite graphs)

Let $G = (V, E)$ be a (general) graph.

- The **(extended) bipartite double** $B_0(G)$ of G has bipartition $V \cup V'$, and edges $\{i, j'\}, \{j, i'\}$ for $\{i, j\} \in E$, and $\{i, i'\}$ for $i \in V$.

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- A pair (A, B) of disjoint subsets of V is biindependent in G
 $\iff (A, B')$ is biindependent in the bipartite graph $B_0(G)$.

Outline

- Motivation
- Relations between the bounds
- Semidefinite and eigenvalue bounds
- Relations between the semidefinite bounds
- Relations to other eigenvalue bounds
- Examples

MOTIVATION

Some applications

- Computing the **maximum number of nodes in a balanced biclique** has applications, e.g., to VLSI design [Al-Yamani et al.'07, Tohhory'06], for analysis of biological data [Yang et al.'05], protein interactions modeling (Mukhopadhyay et al.'14)

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This connection was communicated to us by Frank Vallentin (2020) and motivated our research.

A few details in next slide.

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we get $\varphi(\Gamma) \leq \frac{|\Gamma|}{1+k^{1/3}}$ sharpening Gower's bound $\varphi(\Gamma) \leq \frac{|\Gamma|}{k^{1/3}}$

RELATIONS BETWEEN THE (BALANCED) PARAMETERS

$$\alpha(G) = \max\{|A| + |B| : (A, B) \text{ biindependent in } G\}, \quad \alpha_{\text{bal}}(G)$$

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$$\frac{1}{4}\alpha_{\text{bal}}(G) = \frac{1}{2}\sqrt{g_{\text{bal}}(G)} = h_{\text{bal}}(G) \leq h(G) \leq \frac{1}{2}\sqrt{g(G)} \leq \frac{1}{4}\alpha(G).$$

$$\text{Moreover, } \alpha(G) = \alpha_{\text{bal}}(G) \iff h(G) = \frac{1}{4}\alpha(G) \iff \frac{1}{2}\sqrt{g(G)} = \frac{1}{4}\alpha(G).$$

- The right most inequalities are based on $|A| + |B| \geq 2\sqrt{|A| \cdot |B|}$:

$$h(G) = \frac{|A| \cdot |B|}{|A| + |B|} \leq \frac{1}{2}\sqrt{|A| \cdot |B|} \leq \frac{1}{2}\sqrt{g(G)}$$

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- The non-trivial implication follows from the fact that equality

$$|A| + |B| = 2\sqrt{|A| \cdot |B|} \text{ holds iff } |A| = |B|.$$

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Theorem (L-Polak-Vargas, MOR 2025)

For a bipartite graph G , testing whether $\alpha(G) = \alpha_{\text{bal}}(G)$ is an NP-complete problem.

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Theorem (L-Polak-Vargas, MOR 2025)

For a bipartite graph G , testing whether $\alpha(G) = \alpha_{\text{bal}}(G)$ is an **NP-complete** problem.

Hence, it is **NP-hard** to compute $h(G)$ and $g(G)$.

SEMIDEFINITE AND EIGENVALUE BOUNDS

Semidefinite bounds

$$\boxed{h(G) \leq \frac{1}{2} \sqrt{g(G)} \leq \frac{1}{4} \alpha(G)} \text{ for } G = (V_1 \cup V_2, E) \text{ bipartite graph}$$

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Define the (V_1, V_2) -block matrix $C = \begin{pmatrix} 0 & J \\ J & 0 \end{pmatrix}$, with J all-ones matrix

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$$h_1(G) = \max_{X \in \mathcal{S}^V} \{ \langle C, X \rangle : X \succeq 0, \operatorname{Tr}(X) = 1, X_{ij} = 0 \ (\{i, j\} \in E) \}$$

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Pf: (A, B) biindependent $\leadsto x = \chi^{A \cup B} \leadsto X = xx^\top$

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- $\frac{X}{\operatorname{Tr}(X)}$ is feasible for $h_1(G)$, with value $\frac{|A| \cdot |B|}{|A| + |B|}$

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Compare with **Lovász theta number** (for any G):

$$\vartheta(G) = \max_{X \in S^V} \{ \langle J, X \rangle : X \succeq 0, \operatorname{Tr}(X) = 1, X_{ij} = 0 \ (\{i, j\} \in E) \}$$

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$$\operatorname{las}_1(G) = \max_{X \in S^V} \{ \langle I, X \rangle : \begin{pmatrix} 1 & \operatorname{diag}(X)^T \\ \operatorname{diag}(X) & X \end{pmatrix} \succeq 0, X_{ij} = 0 \text{ } (\{i, j\} \in E) \}$$

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Recall: $\alpha(G) \leq \vartheta(G) = \operatorname{las}_1(G)$, with equality if G is bipartite

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Note: Replace C by J in h_1 (resp., by I in g_1) \leadsto get ϑ (resp., las_1).

Relations between the semidefinite bounds

$$h(G) \leq \frac{1}{2} \sqrt{g(G)} \leq \frac{1}{4} \alpha(G)$$

$$h(G) \leq h_1(G)$$

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Theorem

For any bipartite graph G

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(see later)

About the proofs of

$$\frac{1}{2}\sqrt{g(G)} \leq h_1(G) \leq \frac{1}{2}\sqrt{g_1(G)}$$

- For the inequality $\frac{1}{2}\sqrt{g(G)} \leq h_1(G)$

Use the **dual semidefinite formulation** of $h_1(G)$:

$$h_1(G) = \min_{\lambda \in \mathbb{R}, Z \in S^V} \{ \lambda : \lambda I + Z - C \succeq 0, Z_{ij} = 0 \ (i = j, \{i, j\} \in \overline{E}) \}$$

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and show that $\lambda \geq \frac{1}{2}\sqrt{|A| \cdot |B|}$ if (A, B) is a biindependent pair,
by considering the principal submatrix of $\lambda I + Z - C$ indexed by $A \cup B$.

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- For the inequality $h_1(G) \leq \frac{1}{2}\sqrt{g_1(G)}$, the proof is somewhat analog to the classical proof for showing $\vartheta(G) \leq \text{las}_1(G)$, but more involved

- In comparison, the **dual formulation** of $g_1(G)$ reads:

$$g_1(G) = \min_{\lambda, u \in \mathbb{R}^V, Z \in S^V} \left\{ \lambda : \begin{pmatrix} \lambda & u^\top/2 \\ u/2 & \text{Diag}(u) - C + Z \end{pmatrix} \succeq 0, Z_{ij} = 0 \ (i = j, \{i, j\} \in \overline{E}) \right\}$$

Eigenvalue bounds

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Key proof idea: Use the dual semidefinite programs and restrict to symmetric solutions: $Z = tA_G$, $u = se$

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Lemma (Relation between the eigenvalue bounds)

We have: $h_{\text{eig}}(G) \leq \frac{1}{2} \sqrt{g_{\text{eig}}(G)}$

Relation to Haemers spectral bound

Haemers (2003) studied bicliques and biindependent pairs in general graphs.

Proposition (Link to a bound of Haemers (2003))

Let $G = (V, E)$ a graph, $n = |V|$, and $0 = \mu_1 \leq \mu_2 \leq \dots \leq \mu_n$ the eigenvalues of the Laplacian of G . Let $B_0(G)$ be the **extended bipartite double** of G .

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Moreover, equality $h_1(B_0(G)) = \frac{n}{4} \left(1 - \frac{\mu_2}{\mu_n}\right)$ holds if G is vertex- and edge-transitive.

Recall: $g(B_0(G))$ is the maximum edge-cardinality of a biclique in \overline{G} .

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- However, all natural symmetrizations of these bounds **do not improve** the spectral parameter $h_{\text{eig}}(G)$.

Example: The Hamming cube Q_r

Q_r has vertex set $V = \{0, 1\}^r$, with an edge $\{u, v\}$ if $d_H(u, v) = 1$.

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We show that this is **true** for $r \leq 13$ (using higher order Lasserre bound)

Conclusions

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- The bounds $h_1(G)$ and $g_1(G)$ are the first bounds in a Lasserre-type hierarchy, with finite convergence in $\alpha(G)$ steps.
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- Open problem: Compute $\alpha_{\text{bal}}(Q_r)$ for the Hamming cube, and verify whether it is given by the sequence $a(r-1)$.

THANK YOU !