

# Mixed Integer Programming for Change-Point detection

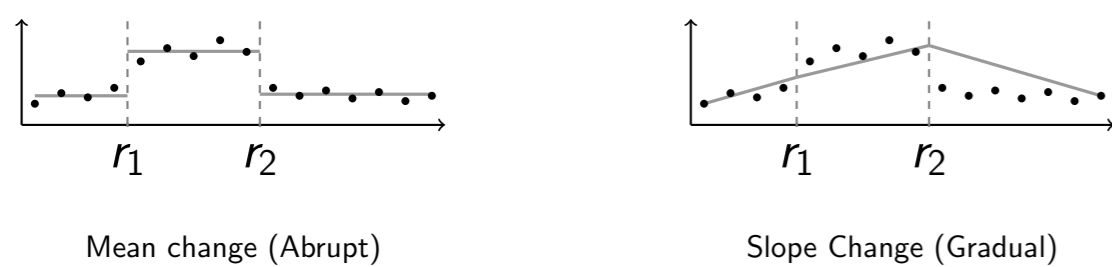
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## 1. Introduction

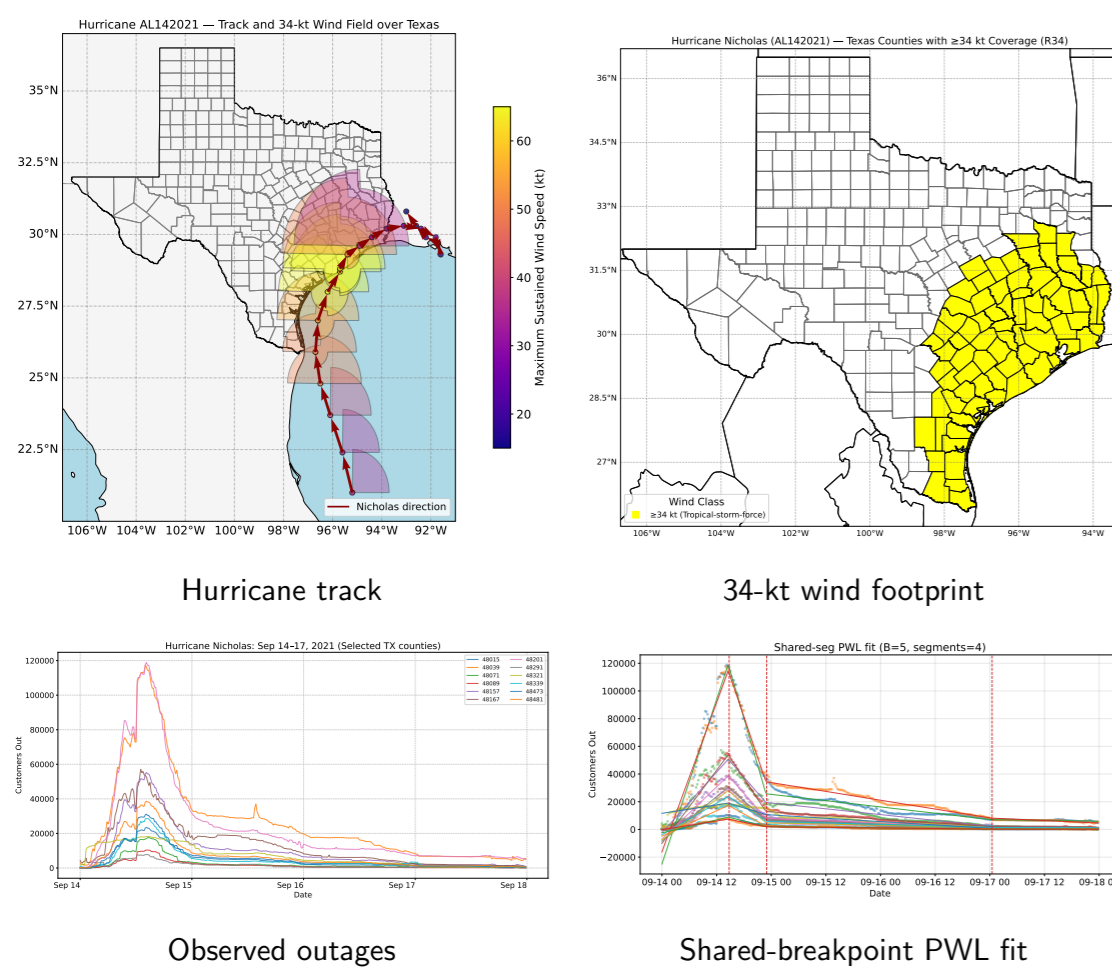
What is Change-point detection?

**Goal:** Locate domain points where the function shifts.



**Applications:** Manufacturing Quality Control (detecting mean shifts), Aircraft Engine Monitoring, Power Systems Transformer Monitoring, and Power Outage Analysis (detecting slope shifts).

**Illustrative example of extreme-weather outage analysis:** Using hurricane trajectory (Hurricane Nicholas, September 2021) and county-level outage data (Texas).



**Interpretation:** Detected change-points can indicate storm impact onset, rapid outage escalation, peak disruption, and recovery initiation.

## 2. Change-point Detection as Contiguous Domain Partitioning

Given observations

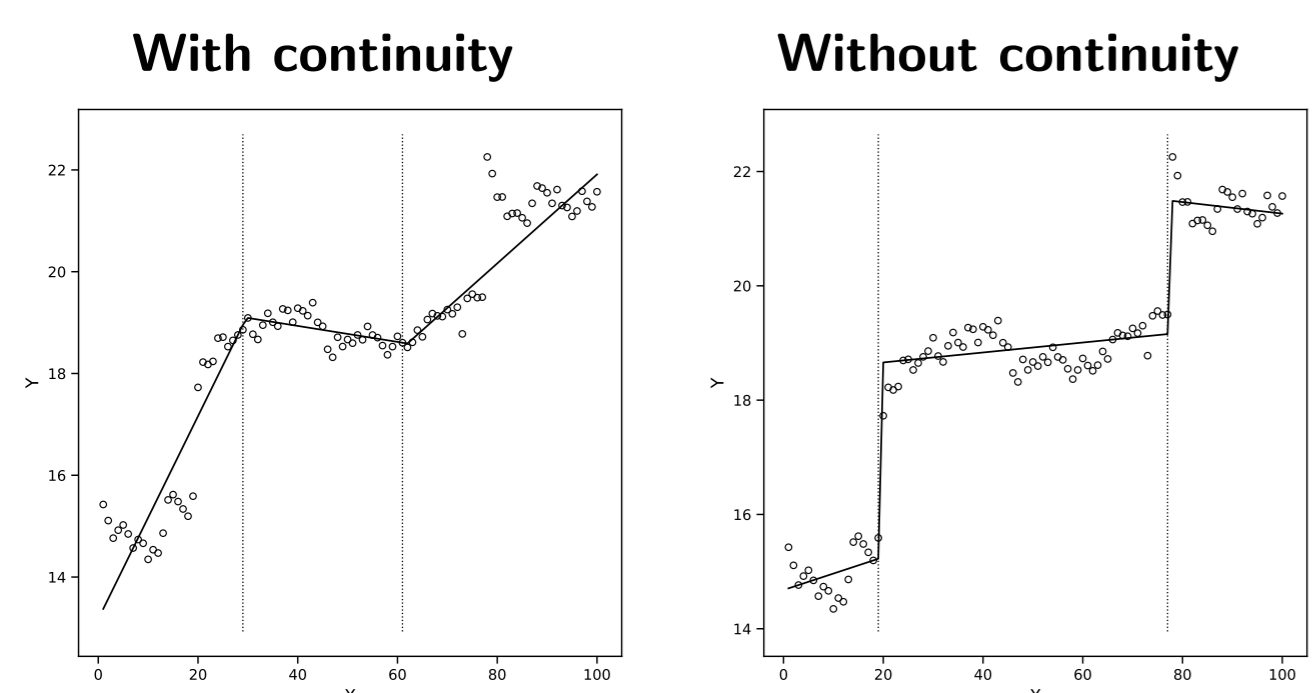
$$(x_t, y_t)_{t=1}^T, \quad x_t, y_t \in \mathbb{R}$$

Partition domain  $[r_0, r_K]$  into  $K$  contiguous segments such that a function piece is only active at a unique partition.

$$f(x) = \begin{cases} f_1(x|\theta_1), & x \in [r_0, r_1], \\ f_2(x|\theta_2), & x \in (r_1, r_2], \\ \vdots \\ f_K(x|\theta_K), & x \in (r_{K-1}, r_K], \end{cases}$$

Decision variables highlighted in red.

## 3. With or Without Continuity



**Contiguous Segmentation Without Continuity:** Polynomial Time - Solvable using Dynamic Programming ✓  
**With Continuity:** NP-Hard - Solvable using MIP - Scalable?

## 4. MIP for Piecewise Linear Fitting

**Objective:**

$$\min \sum_{t=1}^T |y_t - \hat{y}_t|^q, \quad q \in \{1, 2\}$$

**Subject to:**

- **Contiguous Segmentation:** Binary variables  $\delta_{j,t} \in \{0, 1\}$  indicate if point  $t$  is assigned to segment  $j$ .
- **Value fitting:**

$$\hat{y}_t = m_j x_t + c_j \quad \text{if } \delta_{j,t} = 1$$

- **Continuity:** Non-linear, non-convex constraints.

$$m_i r_i + c_i = m_{i+1} r_i + c_{i+1}, \quad \forall i$$

Linearized for piecewise linear fitting by Rebennack and Krasko (2020), *INFORMS Journal on Computing*.

## 5. Motivation and Contributions

**Motivation:** Making MIP-based contiguous segmentation more scalable.

**Key Contributions:** Tighter projection onto relaxed segment assignment variables  $\delta_{j,t}$ ; Faster computation through experiments on stock price datasets from Yahoo Finance.

## 6. Benchmark Formulations

**A. Segment Assignment (Basic):** Goldberg et al. (2021), *Computational Optimization and Applications*. Break-point variables for contiguous segmentation:

**A1. Assignment:** Each observation assigned a segment.

$$\sum_{j=1}^K \delta_{j,t} = 1, \quad \forall t = 1, \dots, T \quad (1)$$

**A2. Breakpoint Localization:**

$$x_t \leq r_j + M_t^{(1)}(1 - \delta_{j,t}) \quad (2)$$

$$x_t \geq r_{j-1} - M_t^{(2)}(1 - \delta_{j,t}) \quad (3)$$

If  $\delta_{j,t}=1$ , then  $x_t \in [r_{j-1}, r_j]$ ; Otherwise, relaxed via big-M.

**B. Segment Assignment (Alternate):** Rebennack & Krasko (2020), *INFORMS Journal on Computing*: Contiguity without breakpoint variables.

**B1. Assignment Constraints:**

$$\sum_{j=1}^K \delta_{j,t} = 1, \quad \forall t = 1, \dots, T \quad (4)$$

**B2. Contiguity Constraints:**

$$\delta_{j+1,t+1} \leq \delta_{j,t} + \delta_{j+1,t} \quad \forall j, t \quad (5)$$

$$\delta_{1,t+1} \leq \delta_{1,t} \quad \forall t \quad (6)$$

$$\delta_{K,t+1} \geq \delta_{K,t} \quad \forall t \quad (7)$$

Enforces contiguous segments, provably by strong induction.

## 7. Proposed Formulation

Narula, Dey & Xie (2026), *arXiv preprint: Contiguous partitioning using nested binary vectors*.

For each  $j = 1, \dots, K-1$ , define

$$X_{j,t} \in \{0, 1\}, \quad t = 1, \dots, T.$$

- $X_{j,t} = 1$ : point  $t$  belongs to one of the first  $j$  segments
- $X_{j,t} = 0$ : point  $t$  belongs to segment  $j+1$  or later

For example:

$$X_{j,\cdot} = [1, 1, 1, 1, 1, 1, 0, 0, 0, 0, 0, 0]$$

The above implies that the first 6 points lie in the first  $j$  segments, so the  $j$ -th breakpoint occurs between  $t = 6$  and  $t = 7$ . To identify the points belonging to segment  $j+1$ , compare two consecutive nested vectors:

$$X_{j,\cdot} = [1, 1, 1, 1, 1, 1, 0, 0, 0, 0, 0, 0]$$

$$X_{j+1,\cdot} = [1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 0, 0, 0]$$

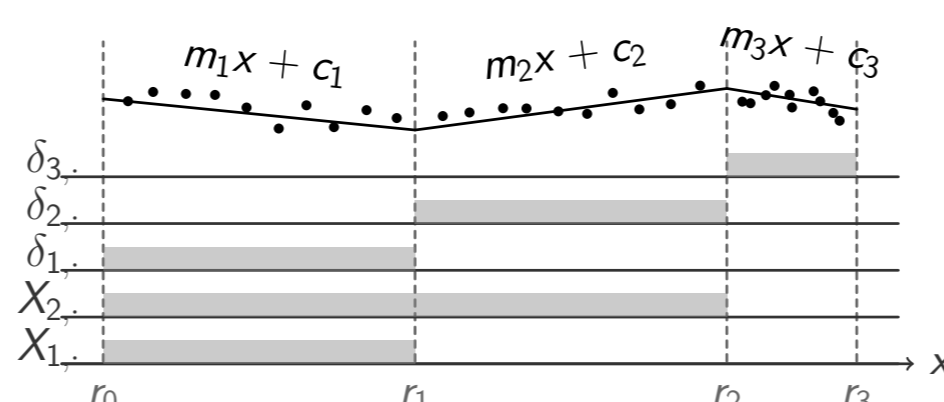
Their difference isolates the points that are in segment  $j+1$ :

$$\delta_{j+1,\cdot} = X_{j+1,\cdot} - X_{j,\cdot}$$

$$\delta_{j+1,\cdot} = [0, 0, 0, 0, 0, 0, 1, 1, 1, 1, 0, 0, 0]$$

Hence, the segment assignments can be recovered from differences of nested binary vectors, as discussed above.

This is also diagrammatically illustrated below:



The above diagram shows how nested binary vectors  $X_{1,\cdot}$  and  $X_{2,\cdot}$  determine the segment-assignment variables  $\delta_{1,\cdot}, \delta_{2,\cdot}, \delta_{3,\cdot}$  in a continuous three-segment piecewise-linear fit.

Hence, the proposed extended formulation is:

**C1. Monotonicity / Nesting Constraints:**

$$\text{Monotonic in Time: } X_{j,t} \geq X_{j,t+1} \quad \forall j, t \quad (8)$$

$$\text{Nested in Segments: } X_{j+1,t} \geq X_{j,t} \quad \forall j, t \quad (9)$$

**C2. Recovering segment assignment:**

$$\delta_{j,t} = \begin{cases} X_{1,t}, & j = 1, \\ X_{j,t} - X_{j-1,t}, & j = 2, \dots, K-1, \\ 1 - X_{K-1,t}, & j = K. \end{cases}$$

## 8. Comparing LP Relaxations

**Prop. 1:** Extended segment assignment polyhedron is integral.

**Proof Sketch:** Constraint matrix is totally unimodular:

- Identity block (from  $\delta$ -definition), appended next to
- Difference constraints (at most two nonzeros:  $+1, -1$ )

**Implication:** Justifies lesser branch and bound effort.

**Prop. 2:** Alternate segment-assignment polyhedron is not integral. **Proof Sketch:** The formulation allows fractional vertices (as shown in the arXiv draft).

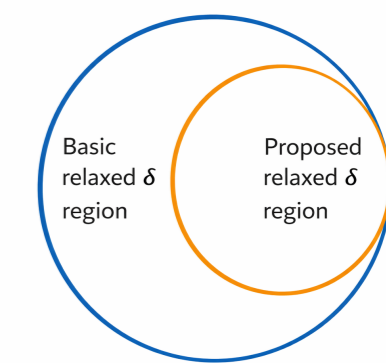
**Prop. 3:** Extended formulation

yields strictly tighter LP

relaxation than Basic

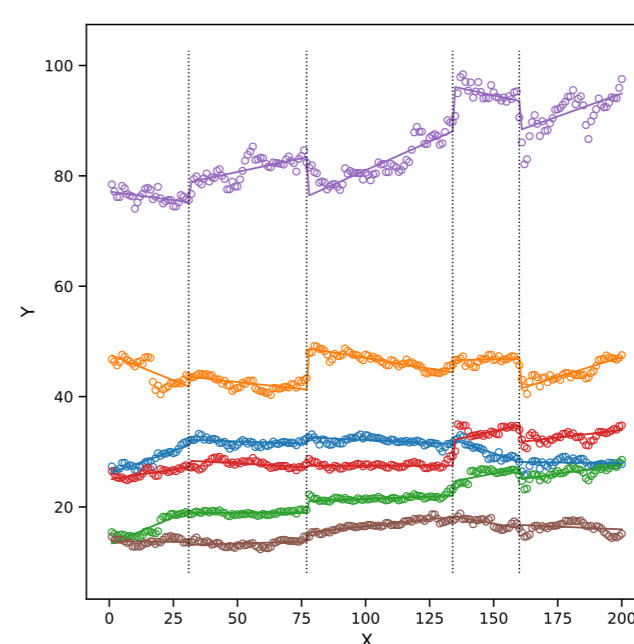
formulation. **Proof Sketch:**

- Basic formulation allows highly fractional, oscillatory relaxed assignments. Proposed formulation enforces total variation bounds on these relaxed assignments.



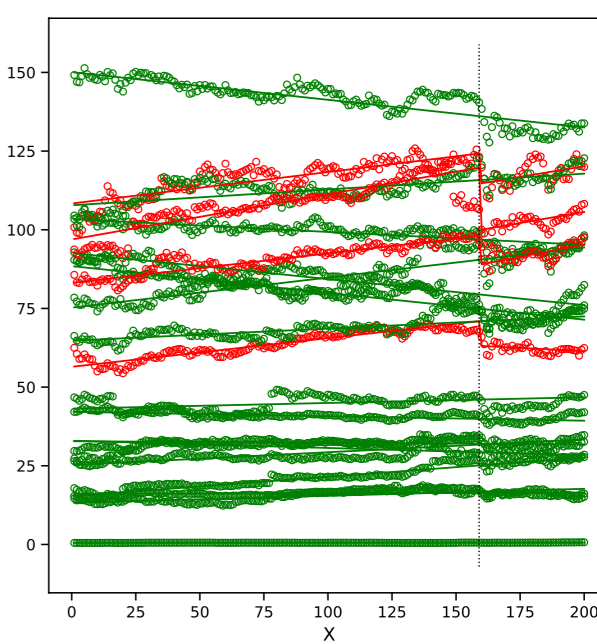
## 9. Multidimensional Extensions

**Multidimensional Change-Point Detection**



$T = 200, K = 5, D = 6$

**Sparse Change-Point Detection**



$D = 20$ , fraction changed = 20%

**Multidimensional:** Fit  $D$  piecewise functions jointly with shared breakpoints, enabling detection of coordinated structural changes across multiple dimensions.

**Sparse Change:** Introduce binary indicators  $\eta_d$  to model whether dimension  $d$  changes, with sparsity enforced through

$$\sum_{d=1}^D \eta_d \leq S.$$

## 10. Computational Comparison

**Experimental Framework:** We compare the Basic, Alternative, and proposed Extended MIP formulations on the following change-point detection tasks:

- Single-dimensional piecewise linear (PWL) fitting: with/without continuity.
- Multi-dimensional PWL fitting,
- Sparse change-point detection,

under both  $\ell_1$  and  $\ell_2$  loss functions. The datasets consist of daily closing stock prices downloaded from Yahoo Finance using Python-based API (yfinance). All formulations were solved using Gurobi 12.0.3 on Apple M3 processor. Performance was evaluated based on the number of instances where a formulation was fastest to converge to globally optimal solution, and total cumulative solver runtime across all instances.

Experiment Setting	Norm	Basic	Alt.	Ext.	Total	
Single-D PWL	No Cont.	$\ell_1$	5	41	54	100
		$\ell_2$	4	27	69	100
	Cont.	$\ell_1$	25	8	67	100
		$\ell_2$	4	36	60	100
Multi-D PWL	$\ell_1$	8	46	46	100	
	$\ell_2$	17	40	43	100	
Sparse CPD	$\ell_1$	2	4	6	12	
	$\ell_2$	7	4	1	12	
<b>Overall</b>			72	206	346	624

(Each entry = number of instances where the formulation was fastest to reach optimality.) **Key Takeaway:** Proposed formulations perform fastest overall.

Experiment Setting	Norm	Basic	Alt.	Ext.	
Single-D PWL	No Cont.	$\ell_1$	3024	2633	3553
		$\ell_2$	1046	661	566
	Cont.	$\ell_1$	41452	35517	24352
		$\ell_2$	48775	23840	21994
Multi-D PWL	$\ell_1$	8112	8834	7760	
	$\ell_2$	3577	3436	3574	
Sparse CPD	$\ell_1$	3046	3883	3359	
	$\ell_2$	9561	10735	10380	
<b>Overall Total</b>		118593	89538	75538	

(Entries = total solver time (seconds) across all instances). **Key Takeaway:** Proposed formulation achieves the lowest cumulative runtime.