

# New Sequence-Independent Lifting Techniques for Cutting Planes and When They Induce Facets

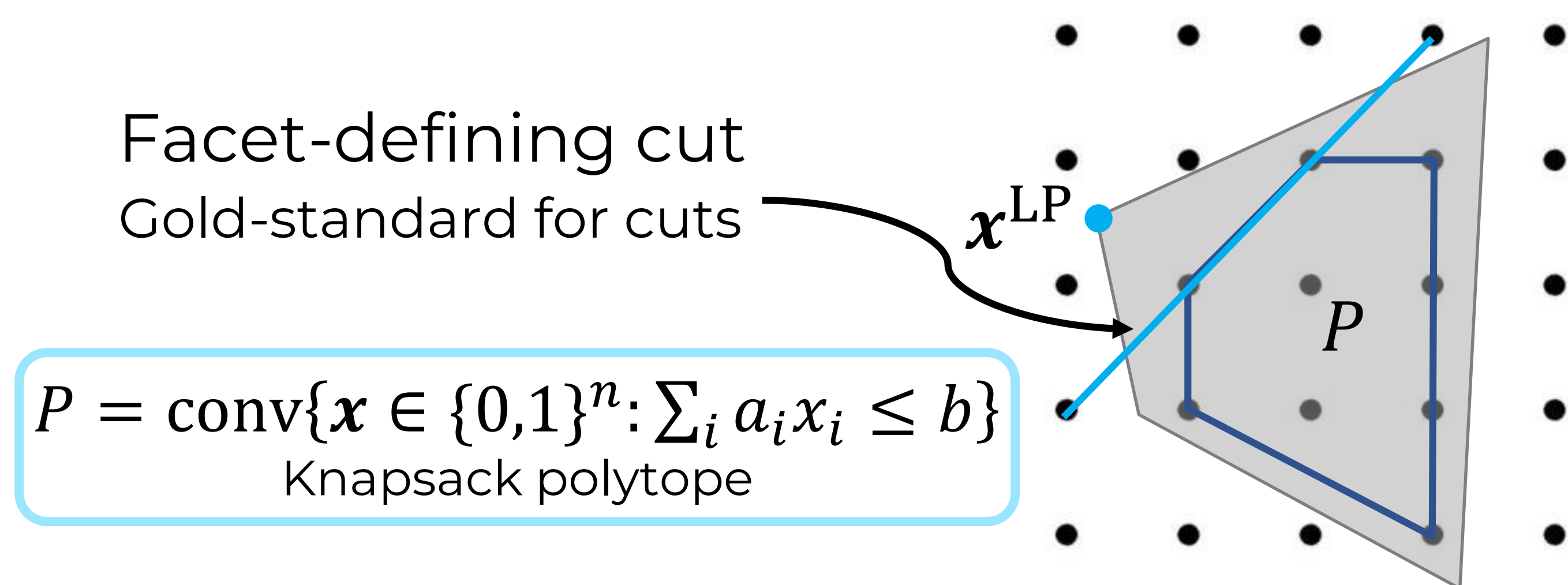
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## Highlights

- Generalization of Gu-Nemhauser-Savelsbergh lifting technique (and correction of their generalization that yields invalid cuts).
- Characterization of when our new lifting function yields facet-defining cuts.
- Experiments with new cover cut families.

## Knapsack Polytopes and Cover Cuts

Facet-defining cut  
Gold-standard for cuts



$$16x_1 + 14x_2 + 13x_3 + 9x_4 + \dots + a_n x_n \leq 44$$

Items 1 – 4 are too heavy  $\rightarrow$  enforce  $x_1 + x_2 + x_3 + x_4 \leq 3$

$C \subseteq \{1, \dots, n\}$  is a cover if  $\sum_{j \in C} a_j > b$

$C$  is a minimal cover if  $\sum_{j \in C \setminus i} a_j \leq b \forall i \in C$

**Minimal cover cut:**  $\sum_{j \in C} x_j \leq |C| - 1$

## Sequence-Independent Lifting

Lifted cover cut:  $\sum_{j \in C} x_j + \sum_{j \notin C} \pi_j x_j \leq |C| - 1$ .

Wolsey (1977): If  $g \leq f$  is superadditive,

$$\sum_{j \in C} x_j + \sum_{j \notin C} g(a_j) x_j \leq |C| - 1$$

is a valid cut;  $g$  is **sequence-independent**.

**Gu-Nemhauser-Savelsbergh (GNS) (2000)**

explicitly constructed a sequence-independent lifting function  $g$ ; very easy to compute (see figure).

$\mu_h$  = weight of  $h$  heaviest items in  $C$ .

$\lambda$  = excess weight of  $C$ .

$\rho_h$  = excess weight of  $C$  if heaviest item replaced with a copy of  $(h + 1)$ st heaviest item.

## Piecewise-Constant (PC) Lifting

**Theorem:** If  $\mu_1 - \lambda \geq \rho_1$ , one can use any slope (uniformly) to get a superadditive lifting fn.

Generalizes GNS lifting and corrects an error in GNS's proposed generalization which yields invalid cuts.

**PC Lifting:** slope-zero piecewise-constant lifting function; all other slopes dominated by PC + GNS.

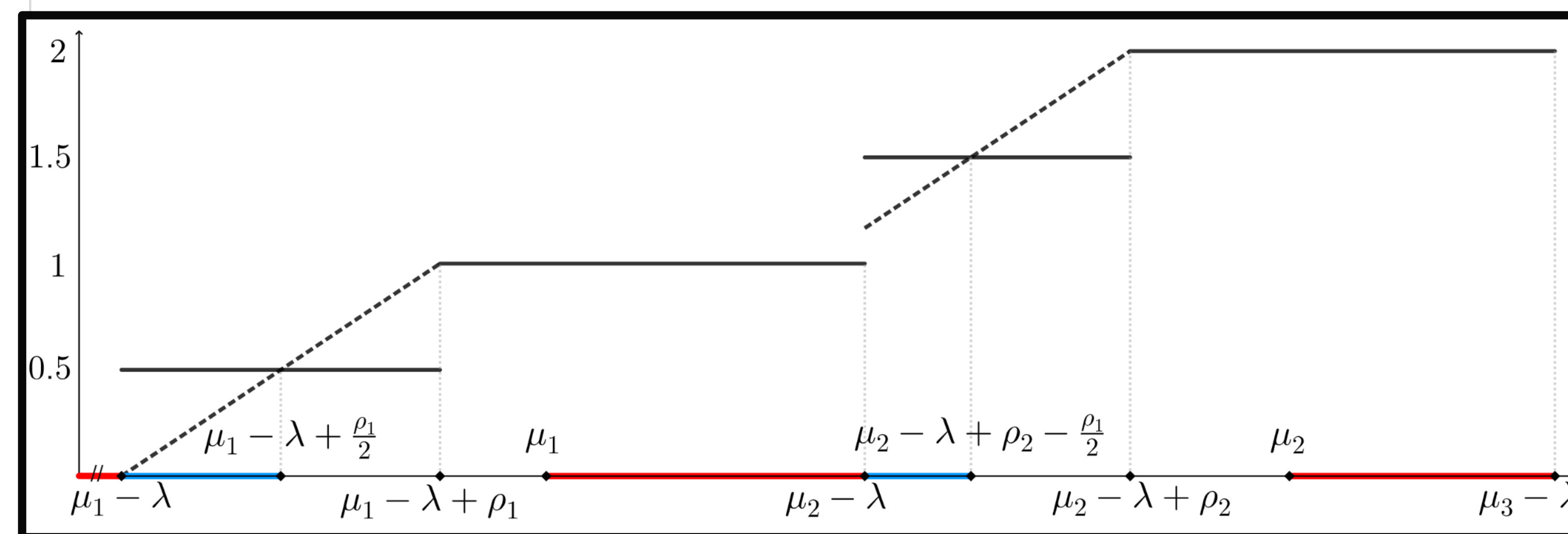


Figure. PC lifting (solid lines) vs GNS lifting (replace with dashed sloped segments).

## Facet-Defining Cuts from PC lifting

**Theorem:** If  $\{a_j : j \notin C\}$  is contained in the **blue** and **red** intervals with at least 3 coeffs. in the leftmost **blue** interval, PC lifting dominates GNS lifting and is facet-defining:

$$\sum_{j \in C} x_j + \sum_{j \notin C} \text{PC}(a_j) x_j \leq |C| - 1 \text{ is a facet of } P.$$

- Proof uses Balas-Zemel (1978) characterization of facets of  $P$  that arise from lifting (BZ does not give a tractable way of deriving such cuts).
- First set of conditions for facet-defining sequence-independent liftings that are efficiently computable from the underlying cover.

$$16x_1 + 14x_2 + 13x_3 + 9x_4 + 9x_5 + 10x_6 + 11x_7 + 23x_8 \leq 44$$

$$\text{GNS: } x_1 + x_2 + x_3 + x_4 + \frac{1}{6}x_5 + \frac{1}{3}x_6 + \frac{1}{2}x_7 + \frac{4}{3}x_8 \leq 3$$

$$\text{PC: } x_1 + x_2 + x_3 + x_4 + \frac{1}{2}(x_5 + x_6 + x_7) + \frac{3}{2}x_8 \leq 3$$

Here, PC strictly dominates GNS and is facet-defining.

## Experimental Evaluation

**Cover cut generation:** cheap, easy to generate families (instead of solving NP-hard separation problem).

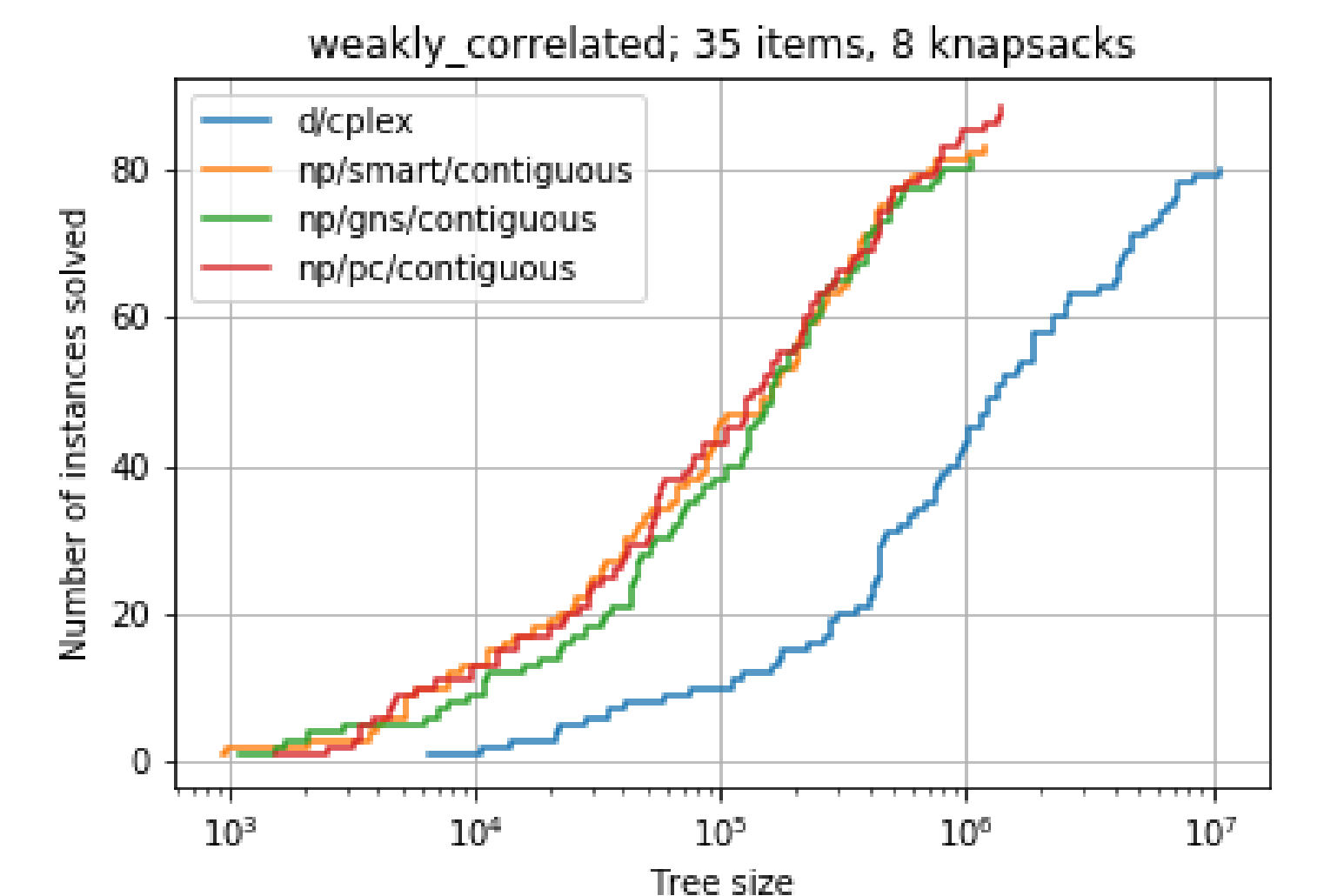
$$10x_1 + 9x_2 + 8x_3 + 7x_4 + 6x_5 + 6x_6 + 5x_7 + 4x_8 \leq 26$$

Contiguous covers  
 $\{1, 2, 3\}, \{2, 3, 4, 5\}, \{3, 4, 5, 6\}, \{4, 5, 6, 7, 8\}$

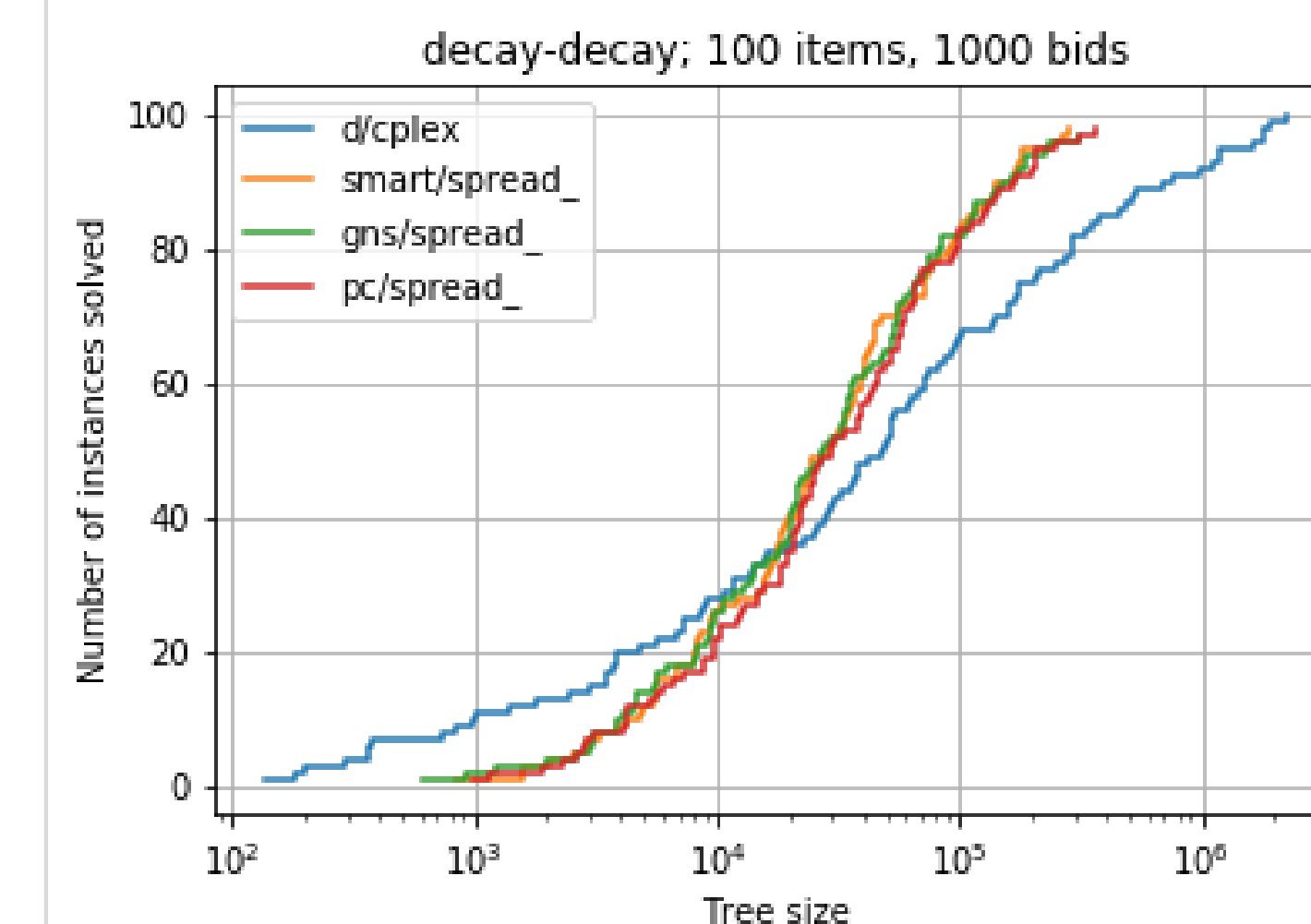
Spread covers  
 $\{1, 5, 6, 7\}, \{2, 5, 6, 7, 8\}, \{3, 5, 6, 7, 8\}, \{4, 5, 6, 7, 8\}$

**Branch-and-Cut integration:** at each node, add the 10 deepest lifted cuts that separate  $\mathbf{x}^{\text{LP}}$ .

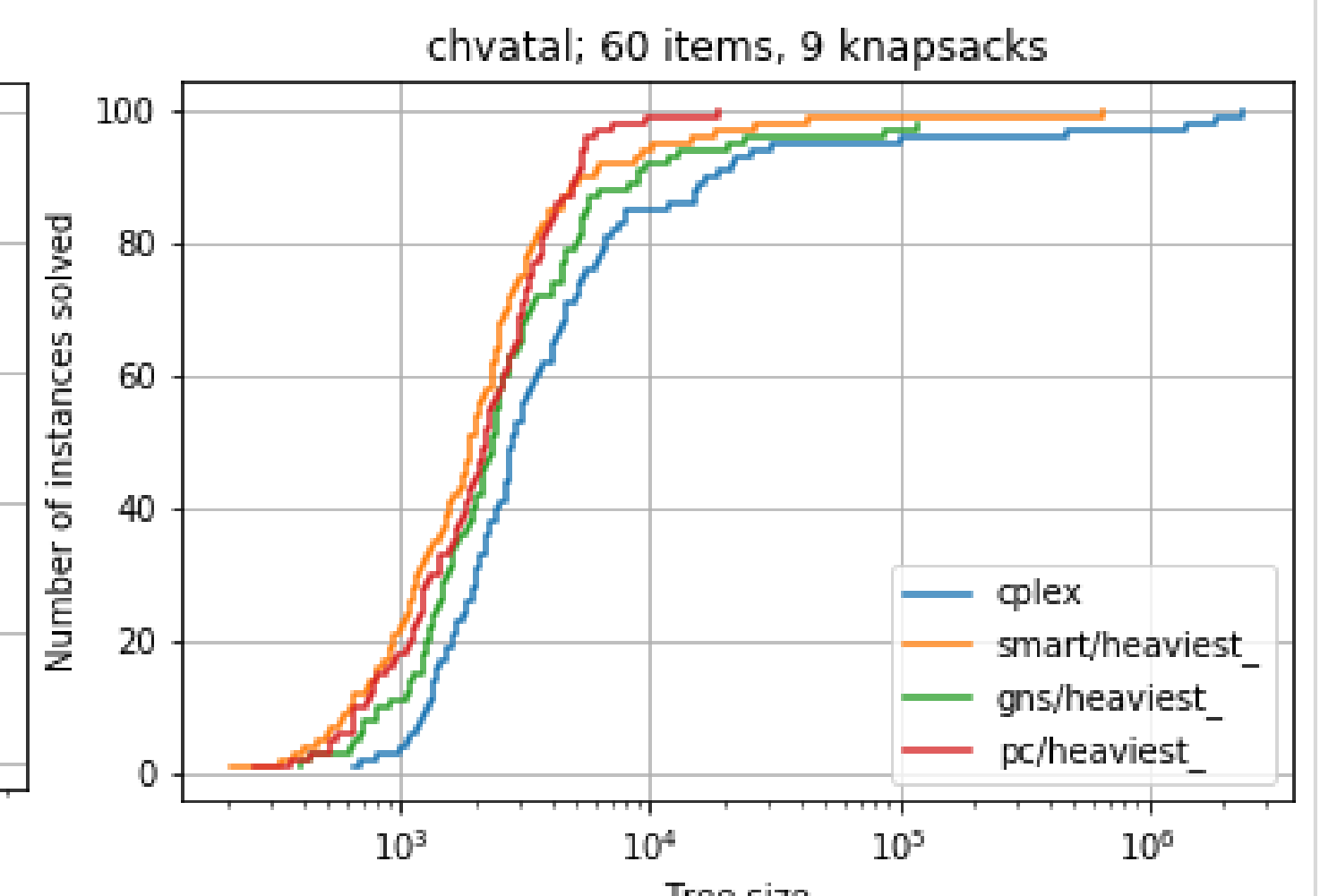
Default CPLEX B&C vs our methods (only presolve off)



Default CPLEX B&C vs our methods (everything off)



Bare-bones CPLEX B&C vs our methods (everything off)



## Practical Issues and Future Research

- Smaller trees didn't translate to run-time improvements, though often we weren't much slower and sometimes we were faster
- More comprehensive suite of experiments needed to see where PC lifting shines.
- Further investigation of desirable numerical properties (half-integral coefficients) of PC lifting.