

# Selecting and Scheduling Cybersecurity Mitigations with Resource Constraints

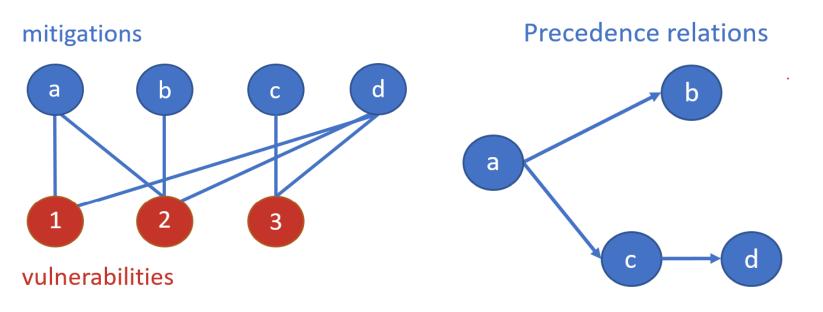
# **Problem Setting**

IT managers regularly face the challenging problem of deploying mitigations to improve cybersecurity. Mitigation selection challenges:

- Appropriately allocating resources over time to achieve security quickly and efficiently.
- Addressing precedence relations when implementing multiple mitigations.
- Prioritizing important vulnerabilities without sacrificing time-efficient overall coverage.

# Problem

What is the best way to schedule the implementation of mitigations subject to resource, budget, and precedence constraints, to achieve maximal coverage of vulnerability nodes when covering a node multiple times gives diminishing returns?



**Existing Model Limitations** 

### • Mitigation selection [2]:

Provide ways to choose mitigations under this multiple-coverage notion, but do not consider deployment in resource-constrained settings.

#### • Resource Constrained Project Scheduling [1]: RCPSPs match our problem's scheduling structure, but cannot model the multiple-coverage objective.

#### IP Model

 $\mathcal{T}$  set  $\{1, ..., T\}$  of time periods  $N \mid \text{set of nodes}$ set of jobs/mitigations set of resources  $R_{-}$ set of precedence relations P $w_{jn}$  benefit of completing  $j \in J$  on  $n \in N$ time required to complete  $j \in J$ cost per period of  $r \in R$  for  $j \in J$  $C_{jr}$ cost of  $j \in J$  for total budget  $C_i$  $b_{rt}$  budget for  $r \in R$  for period  $t \in \mathcal{T}$ total budget (not time indexed) Bpiecewise-linear concave function for objective time-weighting coefficient for objective  $lpha_t$  $z_{nt}$  amount of coverage for  $n \in N$  at time  $t \in \mathcal{T}$ .  $x_{jt} = 1$  if job  $j \in J$  finishes at time  $t \in \mathcal{T}$ , binary  $\max \sum \sum \alpha_t f_n(z_{nt})$  $t=1 n \in N$  $\forall t \in \mathcal{T}, \ n \in N$ s.t.  $z_{nt} \leq \sum \sum w_{jn} x_{js}$  $j \in J s = 1$  $\forall j \in J$  $\sum x_{jt} \le 1$  $\forall j \in J, t = 1, ..., \tau_j - 1$  $x_{jt} = 0$  $t+\tau_j-1$  $\sum_{j \in J} \sum_{s=t}^{J} c_{jr} x_{js} \le b_{rt}$  $\forall t \in \mathcal{T}, \ r \in R$  $\sum_{t=1} \sum_{j \in J} C_j x_{jt} \le B$  $\sum x_{js} \leq \sum x_{is}$  $\forall (i,j) \in P, t = 1, ..., T$  $x_{jt} \in \{0, 1\}$  $\forall j \in J, t \in \mathcal{T}$  $\forall n \in N, t \in \mathcal{T}$  $z_{nt} \ge 0$ 

# Ashley Peper, Jim Luedtke, Laura A. Albert, Department of Industrial and Systems Engineering, University of Wisconsin - Madison

# **Benefit of Integrated Model**

What if we don't combine coverage and scheduling? 

We relax the notion of coverage by removing constraints with z, and using the objective

$$\max \sum_{j \in J} \sum_{t=1}^{T} a_t h_j x_{jt}$$

where  $h_i$  is some estimation of coverage provided by job j, and  $a_t$  is a time-weighting parameter.

#### **2**Select then Schedule (STS)

Ignore scheduling, choose jobs for best coverage. Then use an RCPSP to schedule these jobs.

# **Solving Large Instances**

We propose a rolling horizon heuristic using an interval model derived from [1].

#### Interval Model

- Group time periods into a set of time intervals
- Provides a relaxation of our full model
- Schedules jobs into intervals, which we can use to schedule into time periods (*Int-fast*)

#### Rolling Horizon Heuristic (Int-roll)

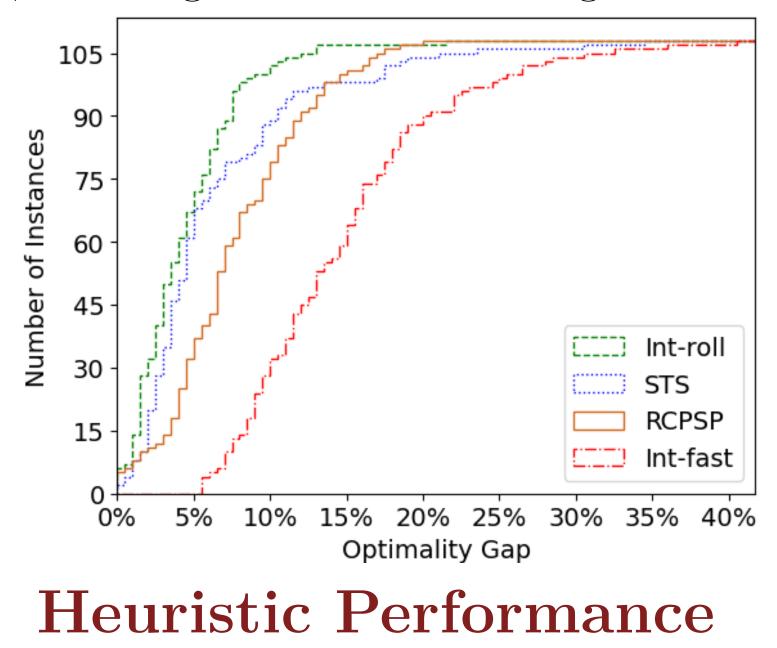
- **1** Use interval model to find a solution.
- Fix some jobs at beginning of horizon.
- <sup>3</sup>Repeat, fixing more jobs later into the horizon at each iteration.

# **Contact Info & Acknowledgements**

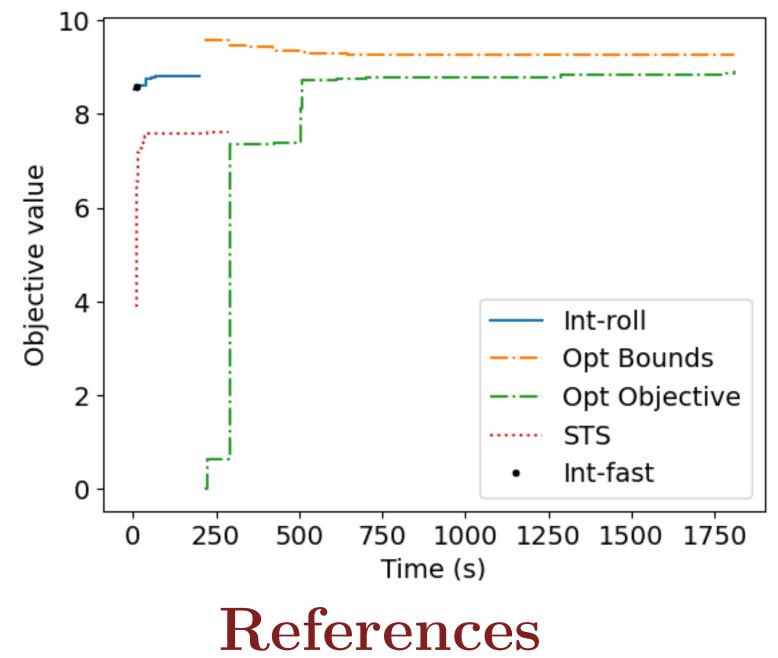
Ashley Peper: apeper2@wisc.edu This work was in part funded by the National Science Foundation Award 2000986.

#### **Computational Comparisons**

• RCPSP and STS average 7% optimality gaps, showing benefit to an integrated model.



• *Int-roll* finds good solutions faster than other methods given a time limit.



- [1] Rodrigo A Carrasco, Diego Fuentes, and Eduardo Moreno. Approximation algorithm for resource- constrained project scheduling problems with net present value objective. arXiv preprint arXiv:2209.02029, 2022.
- [2] Kaiyue Zheng, Laura A Albert, James R Luedtke, and Eli Towle. A budgeted maximum multiple coverage model for cybersecurity planning and management. IISE Transactions, 51(12):1303-1317, 2019.